

Example Write a function **factorial** such that:

If $n \Rightarrow$ a non-negative integer, then $(\text{factorial } n) \Rightarrow n!$.

Recall the following rules for writing recursive functions of 1 argument, which is a proper list or a nonnegative integer:

- When writing a recursive function f , we can first suppose a function f that correctly solves the same problem has already been written.
- Our own version of f can call the supposedly already written f ; but when our version is called with an argument value x , it is only allowed to call the supposedly already written f with an argument value that is valid for f and smaller in size than x .

Assuming **factorial** has already been written correctly, here is a function that works provided $n \Rightarrow$ a nonzero integer:

```
(defun my-factorial (n)
  (let ((X (factorial (- n 1))))
    (* n X)))
```

- We use the fact that:

For example:

$$n * (n-1)! = n!$$

$$5 * 4! = 5 * 4 * 3 * 2 * 1 = 5!$$

Example Write a function **factorial** such that:

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- Our own version of **f** can call the supposedly already written **f**; but when our version is called with an argument value **x**, it is only allowed to call the supposedly already written **f** with an argument value that is valid for **f** and smaller in size than **x**.

Assuming **factorial** has already been written correctly, here is a function that works provided $n \Rightarrow$ a nonzero integer:

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- We use the fact that:

For example:


$$n * (n-1)! = n!$$

$$5 * 4! = 5 * 4 * 3 * 2 * 1 = 5!$$

- Importantly, $n * (n-1)! = n!$ holds even when $n = 1$, as $0! = 1$.
- If $n \Rightarrow 0$, the above definition violates the "call the supposedly already written **f** with an argument value that is valid for **f** and smaller in size" condition, because $(- n 1)$ is not a valid argument value for **factorial** if $n \Rightarrow 0$.

Example Write a function **factorial** such that:

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(defun my-factorial (n)
  (let ((X (factorial (- n 1))))
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```

- If $n \Rightarrow 0$, the above definition violates the "call the supposedly already written *f* with an argument value that is valid for *f* and smaller in size" condition, because $(- n 1)$ is not a valid argument value for **factorial** if $n \Rightarrow 0$.
- To make our function good even when $n \Rightarrow 0$, we add a case:

```
(defun better-my-factorial (n)
  (if (zerop n)
      1
      (let ((X (factorial (- n 1))))
        (* n X))))
```

Example Write a function **factorial** such that:

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(defun better-my-factorial (n)
  (if (zerop n)
      1
      (let ((X (factorial (- n 1))))
        (* n X))))
```

But this still assumes **factorial** has already been written.

Q. How can we write **factorial**?

A. We simply rename **better-my-factorial** to **factorial**!

Example Write a function **factorial** such that:

If $n \Rightarrow$ a non-negative integer, then $(\text{factorial } n) \Rightarrow n!$.

```
(defun factorial factorial (n)
  (if (zerop n) ; base case, where there's no recursive call
      1
      (let ((X (factorial (- n 1))))
        (* n X))))
```

- This definition of **factorial** is not circular, because when **factorial** calls itself it always *passes an argument value that is smaller than the argument value it received.*
- If a recursive call **(factorial (- n 1))** returns the right result, then the call **(factorial n)** returns the right result.
- So, for all positive integers k , if **(factorial i)** returns the right result whenever $i \Rightarrow$ a nonnegative integer $< k$, then **(factorial i)** also returns the right result when $i \Rightarrow k$.
- **Example:** If **(factorial i)** returns the right result when $i \Rightarrow 0, 1, 2$, or 3 , then **(factorial i)** also returns the right result when $i \Rightarrow 4$.

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  (if (zerop n) ; base case, where there's no recursive call
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```

- For all positive integers k , if $(\text{factorial } i)$ returns the right result whenever $i \Rightarrow$ a nonnegative integer $< k$, then $(\text{factorial } i)$ also returns the right result when $i \Rightarrow k$.
 - **Example:** If $(\text{factorial } i)$ returns the right result when $i \Rightarrow 0, 1, 2$, or 3 , then $(\text{factorial } i)$ also returns the right result when $i \Rightarrow 4$.
 - $(\text{factorial } i)$ returns the right result (i.e., 1) when $i \Rightarrow 0$.
- \therefore If $i \Rightarrow$ any nonnegative integer,
then $(\text{factorial } i) \Rightarrow$ the right result (i.e., $i!$).

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- For all positive integers k , if $(\text{factorial } i)$ returns the right result whenever $i \Rightarrow$ a nonnegative integer $< k$, then $(\text{factorial } i)$ also returns the right result when $i \Rightarrow k$.
 - $(\text{factorial } i)$ returns the right result (i.e., 1) when $i \Rightarrow 0$.
- \therefore If $i \Rightarrow$ any nonnegative integer,
then $(\text{factorial } i) \Rightarrow$ the right result (i.e., $i!$).
- Although this function is correct as written, we can improve / simplify the definition by eliminating the LET, because its local variable X is never used more than once.
We then replace the X in $(* n X)$ with $(\text{factorial } (- n 1))$:

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        (* n X))))
```

- Although this function is correct as written, we can *improve / simplify the definition by eliminating the LET, because its local variable X is never used more than once.*

We then replace the `X` in `(* n X)` with `(factorial (- n 1))`:

```
(defun factorial (n)
  (if (zerop n) ; base case, where there's no recursive call
      1
      (* n (factorial (- n 1)))))
```

- As in the case of `length-of`, we've given a written explanation of a possible thought process that leads to this definition, but a Lisp programmer would likely code simple definitions like these without giving any explanation!

- Recursive functions of one argument, which is a list or a nonnegative integer, can often be written in the above way.
- The resulting definition will then have the following form (before possible elimination of the LET):

```
(defun f (e)
  (if (null e)
      value of (f nil)
      (let ((X (f (cdr e))))
        an expression that  $\Rightarrow$  value of (f e)
        and that involves X and, possibly, e ))))
```

OR

```
(defun f (e)
  (if (zerop e)
      value of (f 0)
      (let ((X (f (- e 1))))
        an expression that  $\Rightarrow$  value of (f e)
        and that involves X and, possibly, e ))))
```

- Recursive functions of one argument, which is a list or a nonnegative integer, can often be written as follows:

```
(defun f (e)
  (if (null e) or (zerop e)
      value of (f nil) or (f 0)
      (let ((X (f (cdr e)) or (f (- e 1)))
            an expression that  $\Rightarrow$  value of (f e)
            and that involves X and, possibly, e))))
```

- The `...` expression may have more than one case (as in problem B in Sec. 1 of Lisp Assignment 4): The `...` expression may, e.g., be a **COND** or **IF** expression.
- If there is no case in which **X** is used more than once, then eliminate the LET.
- If the LET isn't eliminated, move any case in which X needn't be used out of the LET. If the LET is eliminated but there's a case where the recursive call's result isn't needed, deal with such cases as base cases--i.e., without making a recursive call.

Example Write a function **evens** such that:

If $L \Rightarrow$ a proper list of integers, then

(evens L) \Rightarrow a list obtained from L by omitting its odd elements.

So (evens '(7 2 -1 4 0 9 2 3)) \Rightarrow (2 4 0 2); (evens nil) \Rightarrow nil.

- Note that the problem specification has this form:

"If $L \Rightarrow$ a proper list of integers, then ..."

This means our function will not be obligated to do anything in particular when its argument value is not a proper list of integers: *It is logically impossible to violate the specification in that case!*

- This is analogous to the meaning of a rule such as:

If you drive on this road, then you must pay a toll.

This rule does not obligate you to do anything if you do not drive on the road in question: *It is logically impossible to violate this rule if you do not drive on the road!*

- If its argument value is not a proper list of integers, then our function **evens** may return any value whatsoever or produce an evaluation error without violating the specification!

Example Write a function `evens` such that:

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So `(evens '(7 2 -1 4 0 9 2 3)) \Rightarrow (2 4 0 2)`; `(evens nil) \Rightarrow nil`.

- If its argument value is **not** a proper list of integers, then our function `evens` may return any value whatsoever or produce an evaluation error without violating the specification!
- The recursive functions you are asked to write will often be specified like this (i.e., with preconditions on argument values that the function may **assume** to be satisfied).
- As a general rule, code that checks that such preconditions are satisfied should **not** be put into short recursive functions: Such checks would complicate/lengthen the code, and may be repeated unnecessarily at every recursive call.
 - Such checks may be done in "gatekeeper" functions that are used by other code to call the recursive functions.
 - Assignments 4 & 5 don't ask you to write such "gatekeeper" functions, but only the recursive functions themselves!

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So `(evens '(7 2 -1 4 0 9 2 3)) \Rightarrow (2 4 0 2)`; `(evens nil) \Rightarrow nil`.

- We'll solve this problem in the way that was described above:

```
(defun f (e)
```

```
  (if (null e)
```

```
    value of (f nil)
```

```
    (let ((x (f (cdr e)))))
```

an expression that \Rightarrow value of `(f e)`
and that involves **x** and, possibly, `e`

Example Write a function **evens** such that:

If $L \Rightarrow$ a proper list of integers, then

$(\text{evens } L) \Rightarrow$ a list obtained from L by omitting its odd elements.

So $(\text{evens } '(7\ 2\ -1\ 4\ 0\ 9\ 2\ 3)) \Rightarrow (2\ 4\ 0\ 2)$; $(\text{evens nil}) \Rightarrow \text{nil}$.

- We'll solve this problem in the way that was described above:

```
(defun evens (L)
```

```
  (if (null L)
```

```
    value of (evens nil)
```

```
    (let ((X (evens (cdr L))))
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      an expression that  $\Rightarrow$  value of (evens L)
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(defun evens (L)
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```

```
    nil
```

```
    (let ((X (evens (cdr L)))))
```

```
      an expression that  $\Rightarrow$  value of (evens L)
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```

- To write the ... expression, let's first consider *one possible value of L , the resulting value of X , and what ...'s value should be for that value of L :*

Suppose $L \Rightarrow (7\ 2\ -1\ 4\ 0\ 9\ 2\ 3)$, so $(\text{cdr } L) \Rightarrow (2\ -1\ 4\ 0\ 9\ 2\ 3)$.

Then $X \Rightarrow (2\ 4\ 0\ 2)$ and ... should $\Rightarrow (2\ 4\ 0\ 2)$.

- For this L , what is a good ... expression? **Ans.: X**

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Then $X \Rightarrow (2\ 4\ 0\ 2)$ and ... should $\Rightarrow (2\ 4\ 0\ 2)$.

- For this L , what is a good ... expression? **Ans.:** X
- Is X a good ... for all non-null values of L ? If not, when is X a good ...? **Ans.** It's good if $(\text{oddp } (\text{car } L))$.

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- We've seen that **X** is a good **...** if $(\text{oddp } (\text{car } L))$. To find a good **...** if $(\text{not } (\text{oddp } (\text{car } L)))$, we try another example:
Suppose $L \Rightarrow (4 \ 2 \ -1 \ 4 \ 0 \ 9 \ 2 \ 3)$, so $(\text{cdr } L) \Rightarrow (2 \ -1 \ 4 \ 0 \ 9 \ 2 \ 3)$.
Then **X** $\Rightarrow (2 \ 4 \ 0 \ 2)$ and **...** should $\Rightarrow (4 \ 2 \ 4 \ 0 \ 2)$.
 - For this L , what is a good **...** expression?
Ans.: $(\text{cons } (\text{car } L) \text{ X})$.
 - Is $(\text{cons } (\text{car } L) \text{ X})$ a good **...** expression for all non-null values of L such that $(\text{not } (\text{oddp } (\text{car } L)))$? Ans. **YES!**

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```

- We've seen that **X** is a good ... if (oddp (car L)).
- We've seen that (cons (car L) **X**) is a good ... if (*not* (oddp (car L))).
- So now we can write ... as:

```
(cond ((oddp (car L)) X)
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- We've seen that **X** is a good ... if (oddp (car L)).
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- So now we can write ... as shown above!

Q. Is there any case in which **X** is used more than once?

A. No! **X** is used just once in each of the 2 cases of the cond.

Example Write a function **evens** such that:

If $L \Rightarrow$ a proper list of integers, then

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Q. Is there any case in which **X** is used more than once?

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- So we can eliminate the LET and substitute (evens (cdr L)) for each occurrence of **X**, to simplify the definition.

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```

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```
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```

```
        (cond ((oddp (car L)) (evens (cdr L)) X)
```

```
              (t (cons (car L) (evens (cdr L)) X))))))
```

- We have **eliminated the LET** and substituted (evens (cdr L)) for each occurrence of **X**, to simplify the definition.
- To further simplify the definition, we can replace (if (null L) nil (cond ...)) with (cond ((null L) nil) ...):

```
(defun evens (L)
```

```
  (cond ((null L) nil)
```

```
        ((oddp (car L)) (evens (cdr L)))
```

```
        (t (cons (car L) (evens (cdr L))))))
```


Recursive Functions of More Than One Argument

- In simple definitions (such as the definitions you are expected to write for Lisp Assignment 4), only one of the arguments of the recursive call needs to have a different value from the corresponding argument of the current call.
- Suppose there are just 2 arguments and the first argument of the recursive call is the argument that has a different value from the corresponding argument of the current call. Then, assuming that argument \Rightarrow a proper list or nonnegative integer, we can often define the function as follows:

```
(defun f (e1 e2)
  (if (null e1) or (zerop e1)
      value of (f nil e2) or (f 0 e2)
      (let ((X (f (cdr e1) e2) or (f (- e1 1) e2) ))
        an expression that  $\Rightarrow$  value of (f e1 e2) and
        that involves X and, possibly, e1 and/or e2 ))))
```

Recursive Functions of More Than One Argument

- In simple definitions (such as the definitions you are expected to write for Lisp Assignment 4), only one of the arguments of the recursive call needs to have a different value from the corresponding argument of the current call.
- Now suppose the second (rather than the first) argument of the recursive call is the argument that has a different value from the corresponding argument of the current call. Then, assuming that argument \Rightarrow a proper list or nonnegative integer, we can often define the function as follows:

```
(defun f (e1 e2)
  (if (null e2) or (zerop e2)
      value of (f e1 nil) or (f e1 0)
      (let ((X (f e1 (cdr e2)) or (f e1 (- e2 1)) ))
        an expression that  $\Rightarrow$  value of (f e1 e2) and
        that involves X and, possibly, e1 and/or e2
      )))
```

Example Without using append, write a function **append-2** such that:

*If L1 \Rightarrow a proper list and L2 \Rightarrow a proper list, then
(append-2 L1 L2) \Rightarrow a list that is equal to (append L1 L2)*

So: (append-2 '(1 2 3 4) '(A B C)) \Rightarrow (1 2 3 4 A B C)

- To solve this problem in the above-mentioned way, we must first decide whether it is the first or the second argument of the recursive call that will have a smaller value than the corresponding argument of the current call.
- Experienced programmers are able to "look ahead" and see which of these two possibilities leads to a good function definition, *but if you can't see which choice is right then just guess*: If your guess doesn't yield a good definition, go back and make the other choice!
- We will attempt to write the function by giving the first argument of the recursive call a smaller value than the corresponding argument of the current call.
- This will turn out to be the right choice!

Example Without using `append`, write a function `append-2` such that:

If $L1 \Rightarrow$ a proper list *and* $L2 \Rightarrow$ a proper list, *then*
 $(\text{append-2 } L1 \ L2) \Rightarrow$ a list that is equal to $(\text{append } L1 \ L2)$

So: $(\text{append-2 } '(1 \ 2 \ 3 \ 4) \ '(A \ B \ C)) \Rightarrow (1 \ 2 \ 3 \ 4 \ A \ B \ C)$

- We will attempt to write the function by giving the first argument of the recursive call a smaller value than the corresponding argument of the current call:

```
(defun append-2 (L1 L2)
```

```
  (if (null L1)
```

```
    value of (append-2 nil L2)
```

What will this be?

```
    (let ((X (append-2 (cdr L1) L2)))
```

an expression that \Rightarrow value of $(\text{append-2 } L1 \ L2)$
and that involves X and, possibly, $L1$ and/or $L2$)))


Example Without using `append`, write a function `append-2` such that:

If $L1 \Rightarrow$ a proper list *and* $L2 \Rightarrow$ a proper list, *then*
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So: $(\text{append-2 } '(1 \ 2 \ 3 \ 4) \ '(A \ B \ C)) \Rightarrow (1 \ 2 \ 3 \ 4 \ A \ B \ C)$

- We will attempt to write the function by giving the first argument of the recursive call a smaller value than the corresponding argument of the current call:

```
(defun append-2 (L1 L2)
  (if (null L1)
      L2
      (let ((X (append-2 (cdr L1) L2)))
        an expression that  $\Rightarrow$  value of  $(\text{append-2 } L1 \ L2)$ 
        and that involves  $X$  and, possibly,  $L1$  and/or  $L2$  )))
```



Example Without using `append`, write a function `append-2` such that:

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So: (append-2 '(1 2 3 4) '(A B C)) \Rightarrow (1 2 3 4 A B C)

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(defun append-2 (L1 L2)
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  (if (null L1)
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```
      (let ((X (append-2 (cdr L1) L2)))
```

an expression that \Rightarrow value of (append-2 L1 L2)
and that involves X and, possibly, L1 and/or L2

)))

- To write the ... expression, let's consider a *possible pair of values of L1 and L2, the resulting value of X, and what ...'s value should be in this case:*
 - Suppose L1 \Rightarrow (1 2 3 4) and L2 \Rightarrow (A B C),
so (cdr L1) \Rightarrow (2 3 4) and X \Rightarrow (2 3 4 A B C).
For this L1 and L2, ... should \Rightarrow (1 2 3 4 A B C).
- Q. What expression (involving X and, possibly, L1 and/or L2) will \Rightarrow (1 2 3 4 A B C)? Ans.: (cons (car L1) X)

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$(\text{append-2 } L1 \ L2) \Rightarrow$ a list that is equal to $(\text{append } L1 \ L2)$

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(defun append-2 (L1 L2)
```

```
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```
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```

```
      (let ((X (append-2 (cdr L1) L2)))
```

an expression that \Rightarrow value of $(\text{append-2 } L1 \ L2)$
and that involves X and, possibly, $L1$ and/or $L2$

```
    )))
```

- Suppose $L1 \Rightarrow (1 \ 2 \ 3 \ 4)$ and $L2 \Rightarrow (A \ B \ C)$,

so $(\text{cdr } L1) \Rightarrow (2 \ 3 \ 4)$ and $X \Rightarrow (2 \ 3 \ 4 \ A \ B \ C)$.

For this $L1$ and $L2$, ... should $\Rightarrow (1 \ 2 \ 3 \ 4 \ A \ B \ C)$.

Q. What expression (involving X and, possibly, $L1$ and/or $L2$)

will $\Rightarrow (1 \ 2 \ 3 \ 4 \ A \ B \ C)$? Ans.: $(\text{cons } (\text{car } L1) \ X)$

Q. Is $(\text{cons } (\text{car } L1) \ X)$ a good ... expression for all

valid values of $L1$ and $L2$ such that $L1 \neq \text{NIL}$?

A. If we're not sure, try *another* pair of values of $L1$ & $L2$.

Example Without using `append`, write a function `append-2` such that:

If $L1 \Rightarrow$ a proper list *and* $L2 \Rightarrow$ a proper list, *then*

$(\text{append-2 } L1 \ L2) \Rightarrow$ a list that is equal to $(\text{append } L1 \ L2)$

```
(defun append-2 (L1 L2)
```

```
  (if (null L1)
```

```
      L2
```

```
      (let ((X (append-2 (cdr L1) L2)))
```

an expression that \Rightarrow value of $(\text{append-2 } L1 \ L2)$
and that involves X and, possibly, $L1$ and/or $L2$

```
    )))
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- Suppose $L1 \Rightarrow (1 \ 2 \ 3 \ 4)$ and $L2 \Rightarrow (A \ B \ C)$,

so $(\text{cdr } L1) \Rightarrow (2 \ 3 \ 4)$ and $X \Rightarrow (2 \ 3 \ 4 \ A \ B \ C)$.

For this $L1$ and $L2$, ... should $\Rightarrow (1 \ 2 \ 3 \ 4 \ A \ B \ C)$.

Q. What expression (involving X and, possibly, $L1$ and/or $L2$)

will $\Rightarrow (1 \ 2 \ 3 \ 4 \ A \ B \ C)$? Ans.: $(\text{cons } (\text{car } L1) \ X)$

- Suppose $L1 \Rightarrow (A \ B \ C \ D \ E \ F)$ and $L2 \Rightarrow (1 \ 2 \ 3 \ 4 \ 5 \ 6 \ 7)$,

so $(\text{cdr } L1) \Rightarrow (B \ C \ D \ E \ F)$ and $X \Rightarrow (B \ C \ D \ E \ F \ 1 \ 2 \ 3 \ 4 \ 5 \ 6 \ 7)$.

For this $L1$ and $L2$, ... should $\Rightarrow (A \ B \ C \ D \ E \ F \ 1 \ 2 \ 3 \ 4 \ 5 \ 6 \ 7)$.

○ $(\text{cons } (\text{car } L1) \ X) \Rightarrow (A \ B \ C \ D \ E \ F \ 1 \ 2 \ 3 \ 4 \ 5 \ 6 \ 7)$ too. Good!

Example Without using `append`, write a function `append-2` such that:

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        an expression that  $\Rightarrow$  value of  $(\text{append-2 } L1 \ L2)$   
        and that involves X and, possibly, L1 and/or L2 )))
```

- Suppose $L1 \Rightarrow (1 \ 2 \ 3 \ 4)$ and $L2 \Rightarrow (A \ B \ C)$,

so $(\text{cdr } L1) \Rightarrow (2 \ 3 \ 4)$ and $X \Rightarrow (2 \ 3 \ 4 \ A \ B \ C)$.

For this $L1$ and $L2$, ... should $\Rightarrow (1 \ 2 \ 3 \ 4 \ A \ B \ C)$.

Q. What expression (involving X and, possibly, $L1$ and/or $L2$)

will $\Rightarrow (1 \ 2 \ 3 \ 4 \ A \ B \ C)$? Ans.: $(\text{cons } (\text{car } L1) \ X)$

- When we are satisfied that $(\text{cons } (\text{car } L1) \ X)$ is a good ... expression for all valid values of $L1$ and $L2$ such that $L1 \neq \text{NIL}$, we complete the above definition!

Example Without using `append`, write a function `append-2` such that:

If $L1 \Rightarrow$ a proper list *and* $L2 \Rightarrow$ a proper list, *then*
 $(\text{append-2 } L1 \ L2) \Rightarrow$ a list that is equal to $(\text{append } L1 \ L2)$

```
(defun append-2 (L1 L2)
  (if (null L1)
      L2
      (let ((X (append-2 (cdr L1) L2)))
        (cons (car L1) X) ))))
```

- Suppose $L1 \Rightarrow (1 \ 2 \ 3 \ 4)$ and $L2 \Rightarrow (A \ B \ C)$,
so $(\text{cdr } L1) \Rightarrow (2 \ 3 \ 4)$ and $X \Rightarrow (2 \ 3 \ 4 \ A \ B \ C)$.
For this $L1$ and $L2$, ... should $\Rightarrow (1 \ 2 \ 3 \ 4 \ A \ B \ C)$.

Q. What expression (involving X and, possibly, $L1$ and/or $L2$)
will $\Rightarrow (1 \ 2 \ 3 \ 4 \ A \ B \ C)$? Ans.: $(\text{cons } (\text{car } L1) \ X)$

Example Without using append, write a function **append-2** such that:

If $L1 \Rightarrow$ a proper list *and* $L2 \Rightarrow$ a proper list, *then*
 $(\text{append-2 } L1 \ L2) \Rightarrow$ a list that is equal to $(\text{append } L1 \ L2)$

```
(defun append-2 (L1 L2)
  (if (null L1)
      L2
      (let ((X (append-2 (cdr L1) L2)))
        (cons (car L1) X))))
```

- **X** is never used more than once, so we *eliminate the LET*:

```
(defun append-2 (L1 L2)
  (if (null L1)
      L2
      (let ((X (append-2 (cdr L1) L2)))
        (cons (car L1) X (append-2 (cdr L1) L2))))))
```

Final version:

```
(defun append-2 (L1 L2)
  (if (null L1)
      L2
      (cons (car L1) (append-2 (cdr L1) L2))))
```